

2-INF-237 Vybrané partie z datových štruktúr

2-INF-237 Selected Topics in Data Structures

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Succinct data structures

We usually count memory in **words** of some size $w \geq \lg n$
each word can hold pointer, index, count, symbol etc.

Now we will count memory in **bits**

Lower bound: to store any $x \in \mathcal{U}$, we need at least $\text{OPT} = \lg |\mathcal{U}|$ bits

Compact data structure uses $O(\text{OPT})$ bits

Succinct data structure uses $\text{OPT} + o(\text{OPT})$ bits

Leading constant 1 plus some lower-order terms

Implicit data structure uses OPT bit plus $O(1)$ words

Uses ordering of elements in an array

Example: binary heap, sorted array

Succinct structure for binary rank and select

Bit vector $A[0..n-1]$

$\text{rank}(i)$ = number of bits set to 1 in $A[0..i]$

$\text{select}(i)$ = position of the i -th bit set to 1

Example:

i	0	1	2	3	4	5	6	7
$A[i]$	0	1	1	0	1	0	0	1

$\text{rank}(3) = 2$, $\text{rank}(4) = 3$

$\text{select}(1) = 1$, $\text{select}(3) = 4$

Goal:

rank , select in $O(1)$ time

structure needs $n + o(n)$ bits of memory

we will concentrate on rank

Succinct structure for rank (Jacobson 1989)

- Divide bit vector to super blocks of size $t_1 = \lg^2 n$
- Divide each super block to blocks of size $t_2 = \frac{1}{2} \lg n$
- Keep rank at each super block boundary
 $O\left(\frac{n}{t_1} \cdot \log n\right) = O(n / \log n) = o(n)$ bits
- Keep rank within super block at each block boundary
 $O\left(\frac{n}{t_2} \cdot \log t_1\right) = O(n \log \log n / \log n) = o(n)$ bits
- Each block stored as a binary number using t_2 bits
 n bits
- For each of 2^{t_2} possible blocks and each query keep the answer
 $O(2^{t_2} \cdot t_2 \cdot \log t_2) = O(\sqrt{n} \log n \log \log n) = o(n)$ bits

Succinct structure for rank (Jacobson 1989)

R1: array of ranks at superblock boundaries

R2: array of ranks at block boundaries within superblocks

R3: precomputed rank for each block type and each position

B: bit array

```
1 rank(i) {
2     superblock = i/t1; //integer division
3     block = i/t2;
4     index = block*t2;
5     type = B[index..index+t2-1];
6     return R1[superblock]+R2[block]+R3[type, i%t2]
7 }
```

Succinct structure for select

- Let $t_1 = \lg n \lg \lg n$, $t_2 = (\lg \lg n)^2$.
- Store $\text{select}(t_1 \cdot i)$ for $i = 0, \dots, n/t_1$;
this divides bit vector into super-blocks of unequal size.
- Large super-blocks of size $\geq t_1^2$: store array of indices of 1 bits.
- Small super-block of size $\leq t_1^2$: repeat with t_1 :
store $\text{select}(t_2 \cdot i)$ within super-block for $i = 0, \dots, n/t_2$;
this divides small super-blocks into blocks of unequal size.
- Large blocks of size $\geq t_2^2$: store relative indices of all 1 bits.
- Small blocks of size $< t_2^2$: store as t_2^2 -bit integer,
plus a lookup table of all answers.

Succinct data structures

- Data structure uses $OPT + o(OPT)$ bits of memory and supports fast operations
- Rank and select on a binary vector of length n in $O(1)$ time

Next:

- Compressed data structures (for rank)
- Wavelet tree for rank over larger alphabet
- Succinct data structure for binary trees

Entropy and compression

Consider alphabet Σ of size σ , probability of $a \in \Sigma$ is p_a

Entropy of this distribution is: $-\sum_{a \in \Sigma} p_a \lg p_a$

Measure of randomness:

Uniform distribution has entropy $\log_2 \sigma$ (max)

If $p_a = 1$ for some $a \in \Sigma$, then entropy 0 (min)

Lossless **compression** of a text consisting of independent identically distributed random symbols with entropy H ,
needs roughly H bits per symbol

Goal: use $-\log_2 p_a$ bit to encode a

Huffman encoding close to that but needs rounding

Arithmetic coding avoids rounding

Compressed structure for rank (Raman, Raman, Rao 2002)

- Compressed size of bit vector + $o(n)$ bits
- Need to reduce the following part:
Each block stored as a binary number using t_2 bits
- Blocks with many 0s or many 1s stored using fewer bits
- For each block store the number of 1s (class)
 $O\left(\frac{n}{t_2} \log t_2\right) = O(n \log \log n / \log n) = o(n)$ bits
- For a block with x 1s store its signature: index in lexicographic order of all binary strings of size t_2 with x 1s
 $\lceil \lg \binom{t_2}{x} \rceil \leq \lg 2^{t_2} = t_2$ bits (overall at most $\frac{n}{t_2} t_2 = n$ bits)
- Rearrange the table with answers for all possible blocks of size t_2
Add signature boundaries in compressed bit vector $o(n)$

Compressed structure for rank (RRR)

$$t_2 = 3$$

Cl	Sig	Length	Block	Answers
0	0	0	000	0 0 0
1	0	2	001	0 0 1
	1		010	0 1 1
	2		100	1 1 1
2	0	2	011	0 1 2
	1		101	1 1 2
	2		110	1 2 2
3	0	0	111	1 2 3

Original bits: 000|101|001|111|111

Number of 1s in each block: 00|10|01|11|11

Index of block: ϵ |01|00| ϵ | ϵ

Where each block starts (within superblock): 0000|0000|0010|0100|0100

Compressed structure for rank (RRR)

rank(i):

- $\text{superblock} = i/t_1$ (integer division)
- $\text{block} = i/t_2$
- $\text{index} = S1[\text{superblock}] + S2[\text{block}]$
- $\text{class} = C[\text{block}]$
- $\text{length} = L[\text{class}]$
- $\text{signature} = B[\text{index}..\text{index} + \text{length} - 1]$
- return $R1[\text{superblock}] + R2[\text{block}] + R3[\text{class}, \text{signature}, i\%t_2]$

Analysis of RRR structure

- Let S be a string in which $a \in \Sigma$ occurs n_a times
- Its **entropy** is $H(S) = \sum_a \frac{n_a}{n} \lg \frac{n}{n_a}$
- RRR structure for bit vector B uses $nH(B) + o(n)$ **bits**

Stirling's approximation of $n!$

$$n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n (1 + O(1/n))$$

$$\ln(n!) = n \ln(n) - n + O(\ln(n))$$

Wavelet tree (Grossi, Gupta, Vitter 2003)

$$\Sigma_0 = \{\$, ., a\} \quad \Sigma_1 = \{e, m, u\}$$

i	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23
S[i]	e	m	a	.	m	a	.	m	a	m	u	.	m	a	m	a	.	m	a	.	e	m	u	\$
B[i]	1	1	0	0	1	0	0	1	0	1	1	0	1	0	1	0	0	1	0	0	1	1	1	0
S0	a.a.a.aa.a.\$					S1																		
						emmmummmemu																		

$$\Sigma_{00} = \{\$, \}, \Sigma_{01} = \{., a\}, \Sigma_{010} = \{.\}, \Sigma_{011} = \{a\}$$

$$\Sigma_{10} = \{e\}, \Sigma_{11} = \{m, u\}, \Sigma_{110} = \{m\}, \Sigma_{111} = \{u\}$$

i	0	1	2	3	4	5	6	7	8	9	10	11	
S0[i]	a	.	a	.	a	.	a	a	.	a	.	\$	
B0[i]	1	1	1	1	1	1	1	1	1	1	1	0	
S00	a.a.a.aa.a.						S01						\$

Store

$$B[i] = 110010010110101001001110$$

$$B0[i] = 111111111110 \quad B1[i] = 011111111011$$

$$B01[i] = 10101011010 \quad B11[i] = 000100001$$

Combination of RRR structure and wavelet trees

- Store binary rank structures in the wavelet tree for text T
overall $nH(T) + o(nH(T))$ bits
- Instead of wavelet tree, store indicator vector for each $a \in \Sigma$
overall $nH(T) + O(n) + o(\sigma n)$ bits
 $O(1)$ per rank query

Dynamic texts (Navarro, Nekrich 2014)

Access, rank, select

Insert/delete character

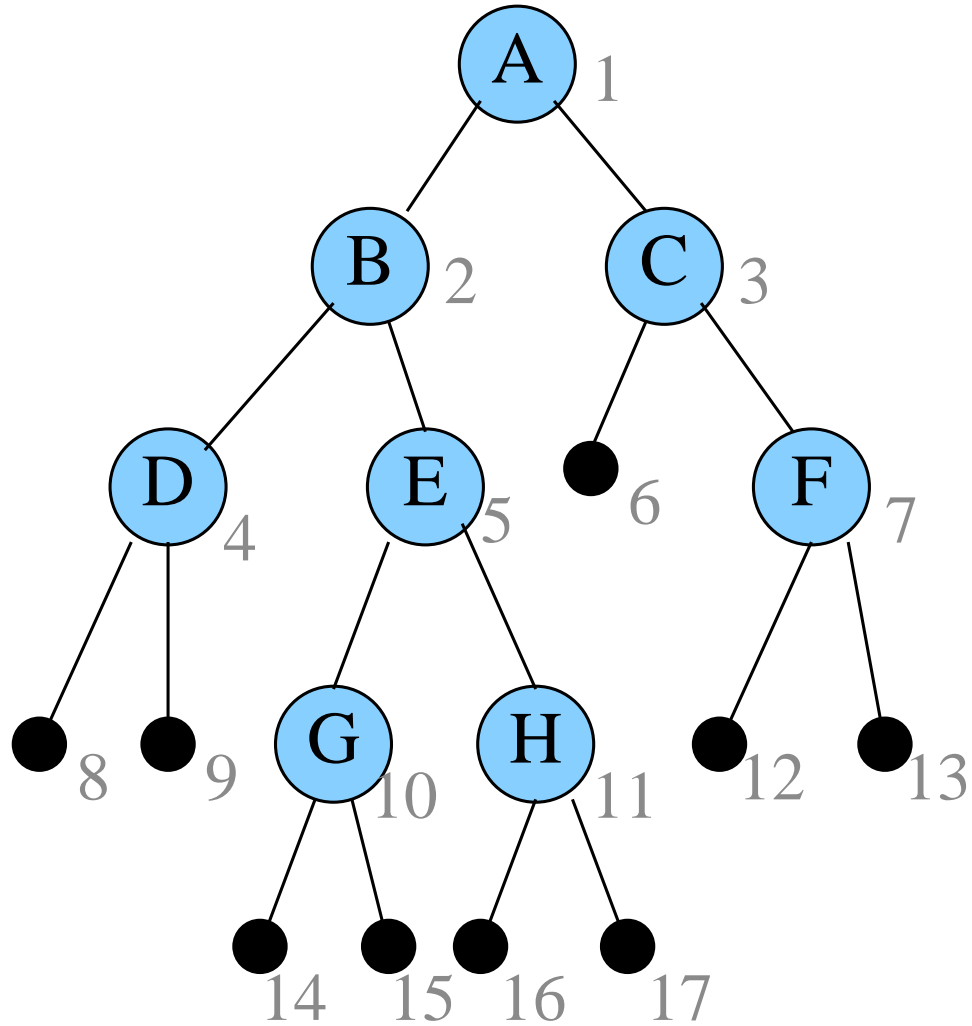
All in $O(\log n / \log \log n)$ amortized

Additional memory $o(n \lg \sigma) + O(\sigma \lg n)$

Succinct binary trees

- Consider all binary trees with n nodes
- Classical trees with pointers use $\Omega(n \log n)$ bits
- OPT is cca $2n$ bits (proof later)
- **Goal:** Use $2n + o(n)$ memory,
support operations left child, right child, parent in $O(1)$
- Add $n + 1$ auxiliary leaves
- Nodes are numbers $\{1, \dots, 2n + 1\}$ in level order (BFS)

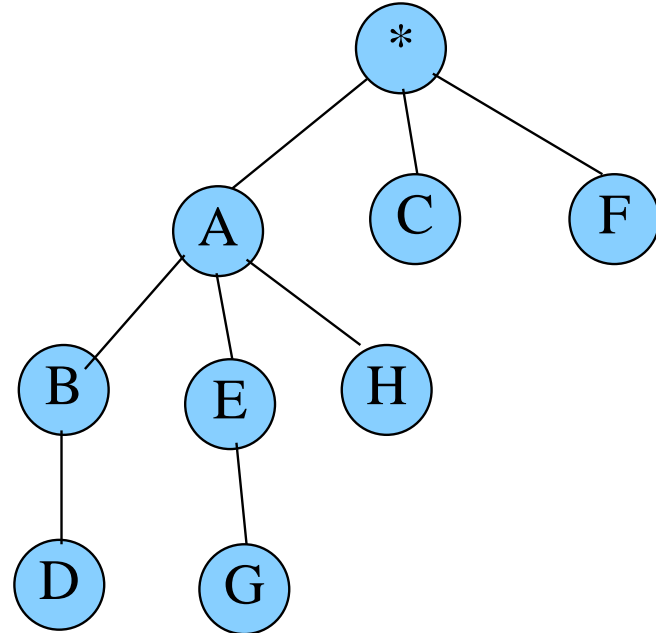
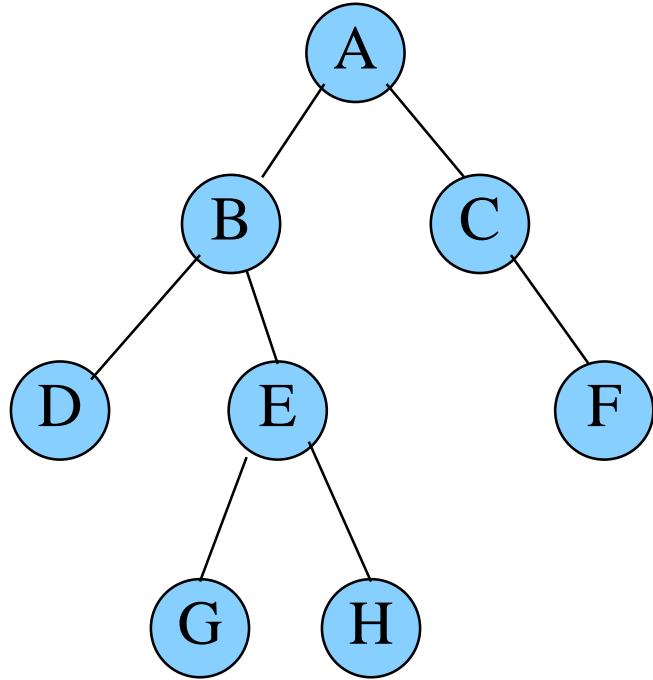
Succinct binary trees: level order representation



Succinct binary trees

- Consider all binary trees with n nodes
- **Goal:** Use $2n + o(n)$ memory,
support operations left child, right child, parent in $O(1)$
- Add $n + 1$ auxiliary leaves
- Nodes are numbers from $\{1, \dots, 2n + 1\}$
Using rank can be mapped to $\{1, \dots, n\}$
These can be then used as indices to arrays with additional data
- Static trees only, construction requires more memory

Equivalence of binary trees and rooted ordered trees



Rooted ordered tree as a well-parenthesized expression:

(((()) (()) ()) () ()

ABDDBEGGEHHACCF

Counting well-parenthesized expressions

$X(n, m, k)$: set of all sequences containing n times 1, m times -1
with all prefix sums $\geq k$

Easy: $|X(n, m, -\infty)|$

Want: $|X(n, n, 0)|$

Prove: $|X(n, n, -\infty) \setminus X(n, n, 0)| = |X(n-1, n+1, -\infty)|$

$$\begin{aligned} \text{Then: } |X(n, n, 0)| &= \binom{2n}{n} - \binom{2n}{n-1} = \frac{(2n)!}{n!n!} - \frac{(2n)!}{(n-1)!(n+1)!} \\ &= \frac{(2n)!(n+1-n)}{n!(n+1)!} = \binom{2n}{n} / (n+1) \end{aligned}$$

Example:

$$|X(3, 3, -\infty) \setminus X(3, 3, 0)| = |X(2, 4, -\infty)| = 15$$

$$|X(3, 3, -\infty)| = 20$$

$$|X(3, 3, 0)| = 20 - 15 = 5 = C_3$$

Counting well-parenthesized expressions

$X(n, m, k)$: set of all sequences containing n times 1, m times -1
with all prefix sums $\geq k$

$$|X(n, n, -\infty) \setminus X(n, n, 0)| = |X(n - 1, n + 1, -\infty)|$$

Consider $n = 3$, first four examples out of 15:

