

Z PREDMÁŠKY

↳ TSP 2-APPROX

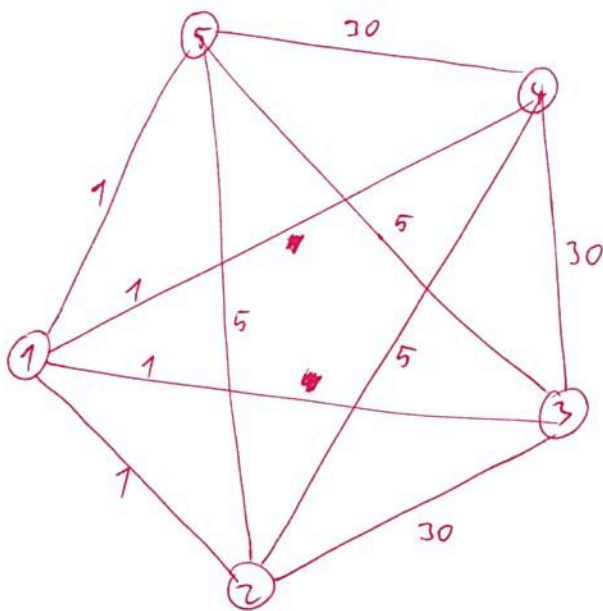
- ↳ $c(T) \leq c(H^*)$ - COST OF MIN. SPANNING TREE \leq OPT. HAMILTONIAN PATH
- ↳ $c(W) = 2 * c(T)$ - FULL WALK W TRAVERSES \forall EDGE TWICE
- ↳ $c(H) \leq c(W)$ - OUR HAMP. PATH FROM REMOVING SOME V FROM W + TRIANGLE INEQUALITY

↓

$$c(H) \leq 2c(H^*)$$

BEZ TRIANGLE INEQUALITY

GRAF S $c(H) \gg 2c(H^*)$



$$c(H) \text{ FROM } 1 = 1-2-3-4-5-1 = 1+30+30+30+1 = 92$$

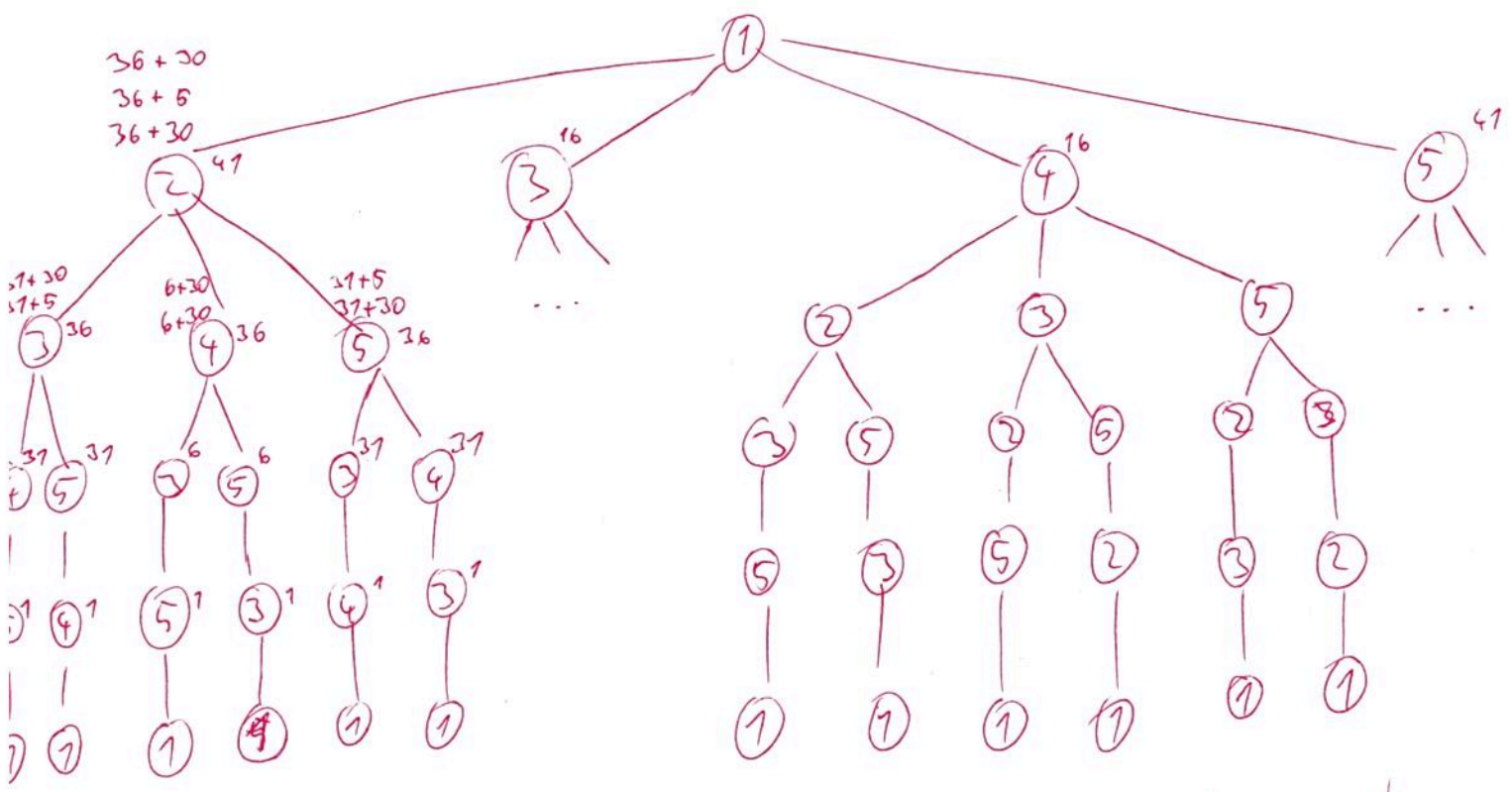
$$c(H) \text{ FROM } 2 = 2-1-3-4-5-2 = 1+1+30+30+5 = 67$$

$$c(H) \text{ FROM } 3 = 3-1-4-5-2-3 = 1+1+30+5+30 = 67$$

$$c(H^*) \text{ FROM } 1 = 1-3-5-2-4-1 = 1+5+5+5+1 = 17$$

- STÁČE BY SME MOGLI TSP RIEŠIŤ !!

↳ DÁ SA ASPOŇ PRE MALÉ n ? \rightarrow BACKTRACK



↳ čas : $O(n!)$
 ↳ DAŤ SA LEPŠIE?

BECCNAK-MELD-MARP

↳ PODPROBLÉM
 ↳ $COST(x_i, S)$
 x_i - VRCHOL GRAFU
 S - SET VRCHOV, KT. OSTÁVAJÚ NA ZVÁŽENIE

D	1	2	3	4	5
1	-	1	1	1	1
2	1	-	30	5	30 ₅
3	1	30	-	30	5
4	1	5	30	-	30
5	1	30 ₅	5	30	-

↳ PODPROBLÉMY - VZŤAHI

$$COST(x_i, S) = \min_{x_j \in S} (COST(x_j, S \setminus \{x_i\}) + D_{ij})$$

↳ BASE CASE

IF $|S|=1$:
 $COST(x_i, S) = D_{i1}$

↳ START

$$COST(x_1, V \setminus \{x_1\})$$

COMPARISON $O(n!)$ vs $O(n^2 \cdot 2^n)$

$$O(100!) \sim O(10^{152})$$

$$O(100^2 \cdot 2^{200}) \sim O(10^{34})$$

TIME COMPLEXITY

$O(2^n)$ ↳ ZVAŽUJEME VŠETKY SUBSETY $V \setminus \{x_1\}$
 ↳ 2^{n-1}

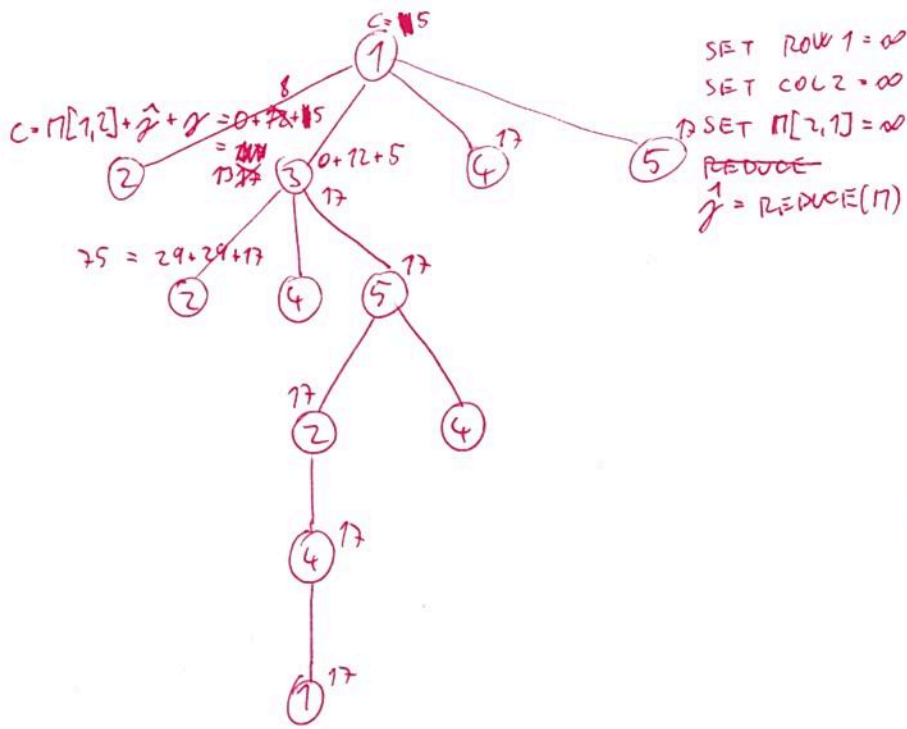
$O(n)$ ↳ PRE KAŽDÝ $\leq n$ SPÚŠŤAME COST FUNKCIU

$O(n)$ ↳ PRE KAŽDÝ COST DERIVUJE ~~MINIMUM~~ MINIMUM Z $\leq n$ HODNÔT

$$\hookrightarrow O(n^2 \cdot 2^n)$$

TSP WITH BRANCH AND BOUND

UB = ∞

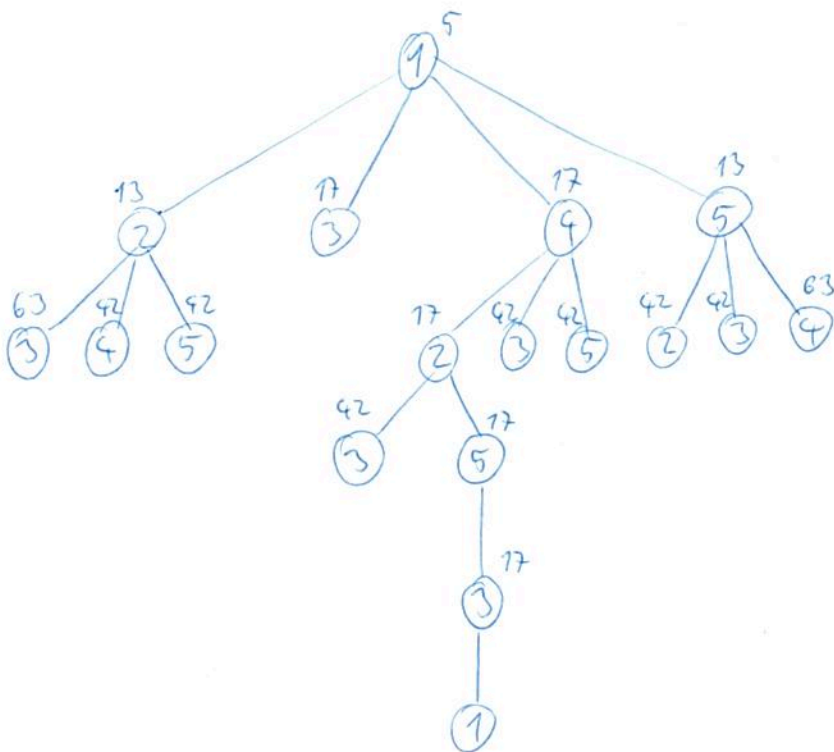


D	1	2	3	4	5
1	-	7	7	1	7
2	7	-	30	5	205
3	7	30	-	30	5
4	1	5	30	-	30
5	7	205	5	30	-

REDUCED D

M	1	2	3	4	5
1	-	0	0	0	0
2	0	-	29	4	29 ⁴
3	0	29	-	29	4
4	0	4	29	-	29
5	0	29 ⁴	4	29	-

$$C = M[z, D_0] + C[z] + \text{COST_OF_REDUCTION}$$



REDUCED COST $z=5$

	1	2	3	4	5
1	-	-	-	-	-
2	-	-	29	4	29 ⁴
3	0	-	-	29	4
4	0	-	29	-	29
5	0	-	4	29	-

0 4 0 0

$z = 8$