

# NEAPPROXIMATE COST VS EOB. TSP

THEOREM (CLRS 35.3):

IF  $P \neq NP$ , THEN FOR ANY CONST.  $\rho \geq 1$ , THERE IS NO POLYNOMIAL TIME APPROX. ALG. WITH APPROX. RATIO  $\rho$  FOR GENERAL TSP.

$$\rho \geq \max\left(\frac{c}{c^*}, \frac{c^*}{c}\right)$$

PROOF: - BY CONTRADICTION

SUPPOSE  $\exists A \in APX$ , WITH  $\rho \geq 1$

USE A TO SOLVE HAMILTONIAN CYCLE PROBLEM IN POLYNOMIAL TIME

HAM  $\in NP$ -COMPLETE

$\Rightarrow P = NP$

LET  $G = (V, E)$  BE INSTANCE OF HAMILTONIAN-CYCLE

LET  $G' = (V, E')$  BE COMPLETE GRAPH ( $E' = \{(u, v) : u, v \in V \wedge u \neq v\}$ )

ASSIGN COST:

$$c(u, v) = \begin{cases} 1 & \text{IF } (u, v) \in E \\ \rho|V| + 1 & \text{ELSE} \end{cases}$$

CONSIDER TSP( $G', c$ )

IF  $G$  CONTAINS HAM. CYCLE  $H \Rightarrow G'$  CONTAINS TOUR OF COST  $|V|$

ELSE  $\Rightarrow \nexists$  TOUR OF  $G'$  CONTAINS  $e \notin E$

$$\hookrightarrow \Rightarrow c(\text{TOUR}) \geq (\rho|V| + 1) + (|V| - 1) = \rho|V| + |V| > \rho|V| \quad (\rho + 1 \text{ FACTOR})$$

BECAUSE  $A \in \rho$ -APX

IF  $G$  HAS HAM  $\Rightarrow A(G', c) \rightarrow H$  WITH COST  $|V|$

IF  $G$  DOES NOT HAVE HAM  $\Rightarrow A(G', c) \nrightarrow H$  WITH COST  $\geq \rho|V| + |V|$

□

GEN. TSP  $\rightarrow$  METRIC TSP

GEN. TSP:

$$D(i, j) \geq 0$$

METRIC TSP:

$$\Gamma = \max_{i, j} D(i, j)$$

$$D'(i, j) = D(i, j) + \Gamma$$

PROOF THAT IT IS METRIC NOW:

$$D'(i, j) + D'(j, k) = D(i, j) + D(j, k) + 2\Gamma \geq 2\Gamma \geq D(i, k) + \Gamma = D'(i, k)$$

PROOF THAT SOLUTION OF METRIC TSP  $\neq$  GEN. TSP:

$$2\text{-APX } A(G', D') \rightarrow \text{TOUR OF COST} \leq 2(C + n\Gamma) = 2C + 2n\Gamma$$

$$2C + 2n\Gamma - n\Gamma = 2C + n\Gamma$$

EACH EDGE  
WAS  $+\Gamma$

$\uparrow$   
THIS IS WAY WORSE THAN 2-APX

# 2-APX METRIC STEINER TREE

$X = R \cup S$ ,  $R$  is required,  $S$  is OPTIONAL

$d: X \times X \rightarrow \mathbb{R}_{\geq 0}$

+ TRIANGLE INEQUALITY

FIND TREE  $T = (V, E)$ , WHERE  $V$  IS ANY SET  $R \subseteq V \subseteq X$  S.T.

$COST(T) = \sum d(u, v)$  IS MINIMIZED

1) COMPLETELY DISREGARD OPTIONAL VERTICES  $S$

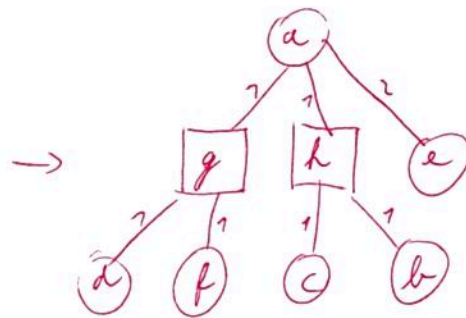
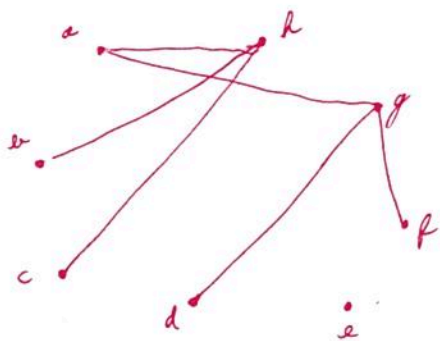
2) FIND MIN. SPAN. TREE

PROOF:  $(X = (R \cup S), d)$  - INSTANCE OF METRIC STEINER TREE

$T = (V, E)$

$T' = (R, E')$  S.T.

$COST(T') \leq 2 COST(T)$



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OSTATNŮ 2

$c=16 \Rightarrow a \rightarrow g \rightarrow d \rightarrow g \rightarrow f \rightarrow g \rightarrow h \rightarrow c \rightarrow h \rightarrow b \rightarrow h \rightarrow a \rightarrow e \rightarrow$

↓  
x OPTIONAL  
x USED

$c=10 \Rightarrow a \rightarrow d \rightarrow f \rightarrow c \rightarrow b \rightarrow e$